

Symmetric distinguishability as a quantum resource

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Joint work with: Nilanjana Datta, Gilad Gour, Xin Wang and
Mark M. Wilde

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Quantum resources:

From mathematical foundations to operational characterisation
Singapore, 6 December 2022

Introduction

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¹X. Wang , M.M. Wilde, Phys. Rev. Res., 1(3):033170, 2019.

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- Fundamental objects involved are **elementary quantum sources**. We study transformations between named sources via **free operations**.

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- Here we develop the **resource theory of symmetric distinguishability (RTSD)** (Earlier work: resource theory of asymmetric distinguishability¹).
- Fundamental objects involved are **elementary quantum sources**. We study transformations between named sources via **free operations**.
- Key result: The **quantum Chernoff divergence** is fundamental exchange rate in RTSD.

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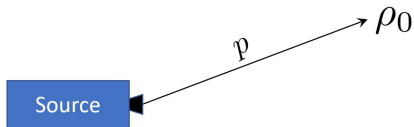
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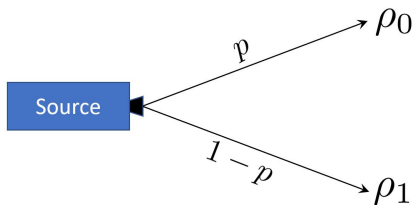
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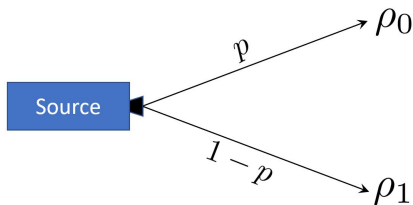
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Such sources can be represented by classical quantum (c-q) state

$$\rho_{XA} := p|0\rangle\langle 0| \otimes \rho_0 + (1 - p)|1\rangle\langle 1| \otimes \rho_1$$

Symmetric Hypothesis testing

Task: Given quantum source $\rho_{XA} \equiv (p, \rho_0, \rho_1)$, discriminate between the states ρ_0 and ρ_1 .

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Infinite resource: ρ_0 and ρ_1 perfectly distinguishable ($\rho_0 \perp \rho_1$) and hence $\text{SD}(\rho_{XA}) = \infty$.

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Symmetric Hypothesis testing: Asymptotic i.i.d. setting

Source emitting n copies of quantum states represented by

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Hence, $\rho_{XA}^{(n)}$ becomes an infinite resource as $n \rightarrow \infty$.

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Resource theory of symmetric distinguishability

Free operations for transformation between quantum sources

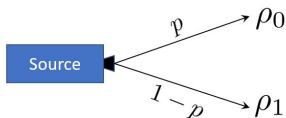
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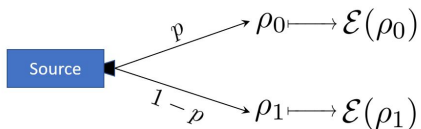


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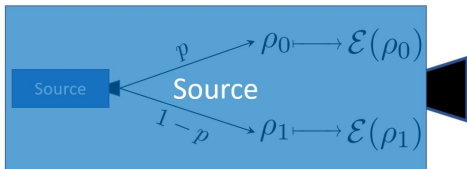


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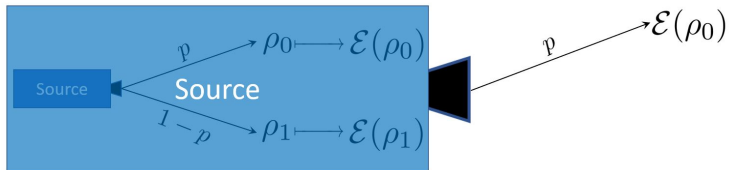


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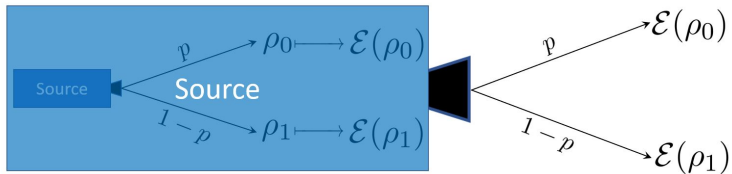


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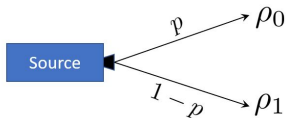
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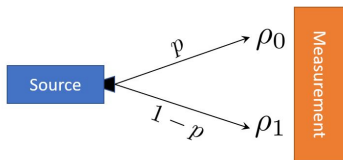


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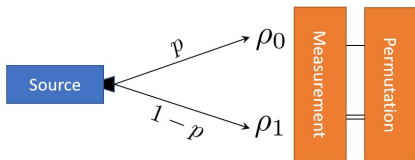


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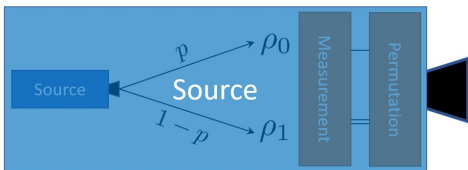


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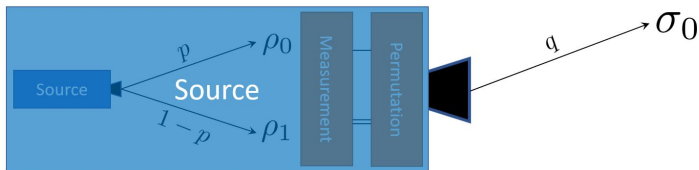


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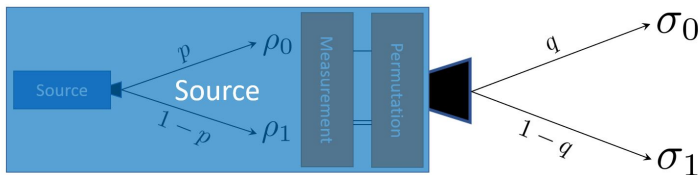


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Monotonicity of minimum error probability:

$$p_{\text{err}}(\mathcal{N}(\rho_{XA})) \geq p_{\text{err}}(\rho_{XA}), \quad \text{for all } \mathcal{N} \in \text{CDS}.$$

Allowing for errors in the transformations

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Error measure:

We allow for errors in the transformations which will be measured by the **scaled trace distance**

$$D'(\tilde{\sigma}_{XB}, \sigma_{XB}) := \frac{\|\tilde{\sigma}_{XB} - \sigma_{XB}\|_1}{2p_{\text{err}}(\sigma_{XB})}.$$

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Important properties:

- **Monotonic** under free operations:

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- **Key inequality** for proving converses in RTSD:

$$p_{\text{err}}(\tilde{\sigma}_{XB}) \leq (D'(\tilde{\sigma}_{XB}, \sigma_{XB}) + 1)p_{\text{err}}(\sigma_{XB}).$$

Asymptotic transformations

What is the optimal (i.e. largest) achievable rate $\frac{m}{n}$ such that

$$\rho_{XA}^{(n)} \equiv (p, \rho_0^{\otimes n}, \rho_1^{\otimes n}) \longmapsto \sigma_{XB}^{(m)} \equiv (q, \sigma_0^{\otimes m}, \sigma_1^{\otimes m}) \quad (1)$$

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We will denote by

$$R(\rho_{XA} \mapsto \sigma_{XB}), \quad \text{under free operations being CPTP}_A,$$
$$R^*(\rho_{XA} \mapsto \sigma_{XB}), \quad \text{under free operations being CDS},$$

the **optimal asymptotic rate** of the transformation (1) such that the D' -error vanishes as $n \rightarrow \infty$.

Theorem

(Optimal asymptotic rate of transformations in RTSD)

Under free operations being CDS :

$$R^*(\rho_{XA} \mapsto \sigma_{XB}) = \frac{\xi(\rho_0, \rho_1)}{\xi(\sigma_0, \sigma_1)}.$$

Under free operations being CPTP_A and $p = q$:

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- The result shows that the RTSD is asymptotically reversible.
- The strong converse property holds: Any transformation with rate larger than $\frac{\xi(\rho_0, \rho_1)}{\xi(\sigma_0, \sigma_1)}$ leads to infinite D' -error as $n \rightarrow \infty$.

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Achievability: We define task of **SD-distillation and dilution** with M -golden unit $\gamma_{XQ}^{(M)}$:

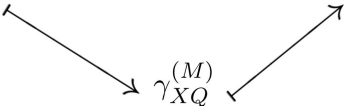
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The diagram illustrates the relationship between the input state $\rho_{XA}^{(n)}$, the output state $\sigma_{XB}^{(m)}$, and the M -golden unit $\gamma_{XQ}^{(M)}$. The top part shows a mapping from $\rho_{XA}^{(n)} \equiv (p, \rho_0^{\otimes n}, \rho_1^{\otimes n})$ to $\sigma_{XB}^{(m)} \equiv (q, \sigma_0^{\otimes m}, \sigma_1^{\otimes m})$. Below this, two arrows point from $\rho_{XA}^{(n)}$ and $\sigma_{XB}^{(m)}$ towards the central node $\gamma_{XQ}^{(M)}$, indicating that both states are related to or derived from this unit.

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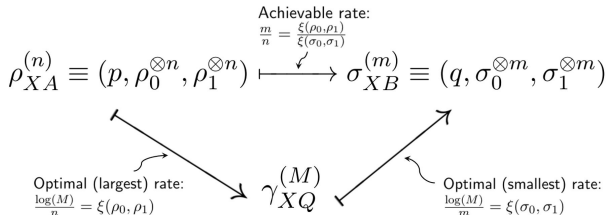
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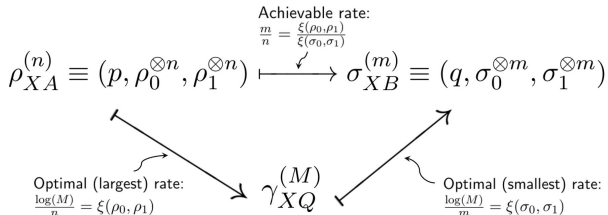
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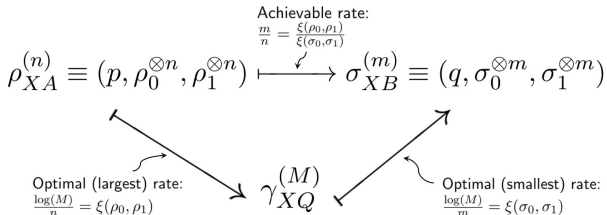
Converse: Assume there exists free operation \mathcal{N} s.t.

$\mathcal{N}(\rho_{XA}^{(n)}) \approx_{D'}^{\varepsilon} \sigma_{XB}^{(m)}$. Then

$$p_{\text{err}}(\rho_{XA}^{(n)}) \leq p_{\text{err}}(\mathcal{N}(\rho_{XA}^{(n)}))$$

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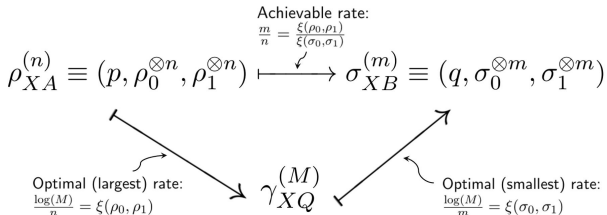
Converse: Assume there exists free operation \mathcal{N} s.t.

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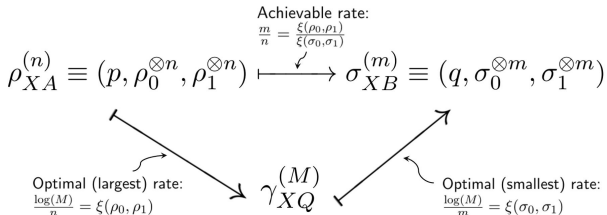
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 \implies \frac{\xi(\rho_0, \rho_1)}{\xi(\sigma_0, \sigma_1)} &\geq \frac{m}{n} + o(1) \quad \text{as } n \rightarrow \infty.
 \end{aligned}$$

Definition (Golden unit)

For $M \in [1, \infty]$ and $q \in (0, 1)$ consider the (M, q) -golden unit

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For M large enough, $\gamma_{XQ}^{(M,q)}$ has well-behaved SD as

$$p_{\text{err}} \left(\gamma_{XQ}^{(M,q)} \right) = \frac{1}{2M} \quad \text{and hence} \quad \text{SD} \left(\gamma_{XQ}^{(M,q)} \right) = \log(M).$$

SD-distillation

Distil golden unit from initial source ρ_{XA} via free operations:

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What is the largest possible M which can be distilled from ρ_{XA} ?

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One-shot exact distillable SD under free operations being CPTP_A

$$\xi_d(\rho_{XA}) := \log \left(\sup \left\{ M \mid \rho_{XA} \xrightarrow{\text{CPTP}_A} \gamma_{XQ}^{(M,p)} \right\} \right)$$

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Allowing for D' -error $\varepsilon > 0$ in the above transformations, the **one-shot approximate distillable SD** will be denoted by $\xi_d^\varepsilon(\rho_{XA})$ (CPTP_A) and $\xi_d^{*,\varepsilon}(\rho_{XA})$ (CDS).

Asymptotic SD-distillation

Denoting again $\rho_{XA}^{(n)} \equiv (p, \rho_0^{\otimes n}, \rho_1^{\otimes n})$ we have:

Theorem

Optimal asymptotic rates of exact and approximate SD-distillation (for all $\varepsilon > 0$) are given by quantum Chernoff divergence:

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Dilute target source ρ_{XA} from golden unit via free operations:

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Here, $d_T(\rho_0, \rho_1) := \max \left\{ D_{\max}(\rho_0 \parallel \rho_1), D_{\max}(\rho_1 \parallel \rho_0) \right\}$,⁶ and $D_{\max}(\rho_0 \parallel \rho_1) := \inf \left\{ \lambda \mid \rho_0 \leq 2^\lambda \rho_1 \right\}$.

⁶A.C. Thompson, Proc. Amer. Math. Soc., 14(3):438–443, 1963.

⁷N. Datta, IEEE Trans. on Inf. Theo., 55(6):2816–2826, 2009.

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Thanks for your attention!

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